

# Introduction to Communicating Sequential Process (CSP) (Lecture 2)

Mannheim, September 2007

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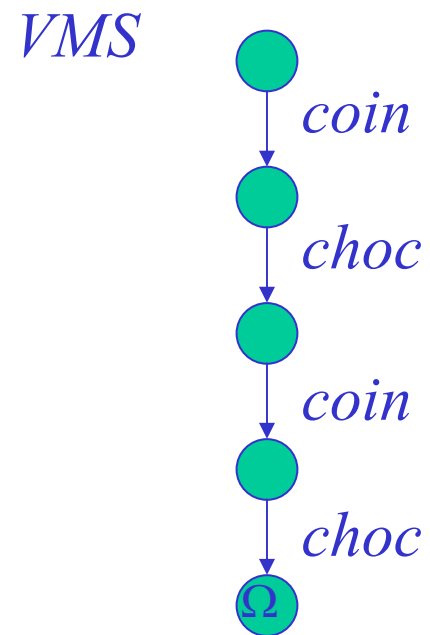
- Pictures
- Basic constructs laws
- Choices
  - Internal (nondeterministic) choice
  - External (deterministic) choice
  - Conditional choice

# Pictures

- It may be useful to make a pictorial representation of a process as a tree structure, where
  - nodes are states
  - arrows are transitions between states
- The tree includes a starting state

# Pictures

$VMS = (\textit{coin} \rightarrow \textit{choc} \rightarrow \textit{coin} \rightarrow \textit{choc} \rightarrow \textit{STOP})$

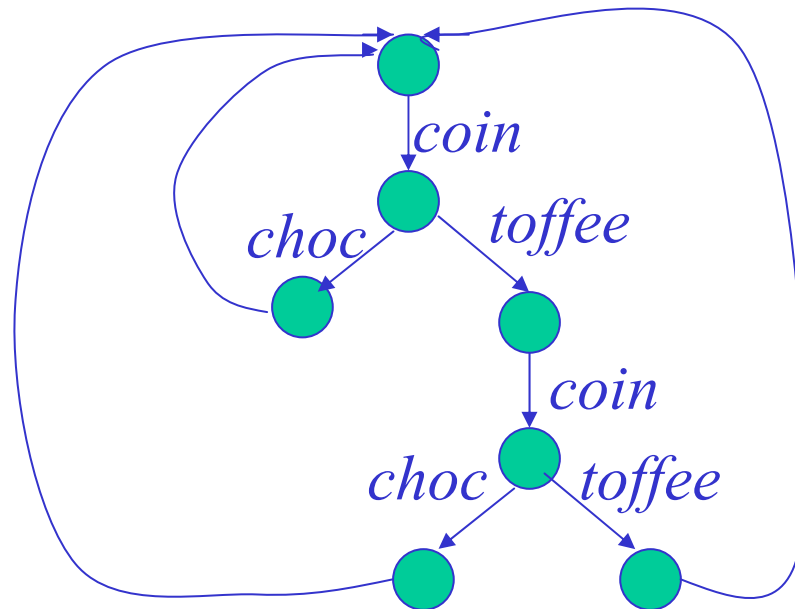
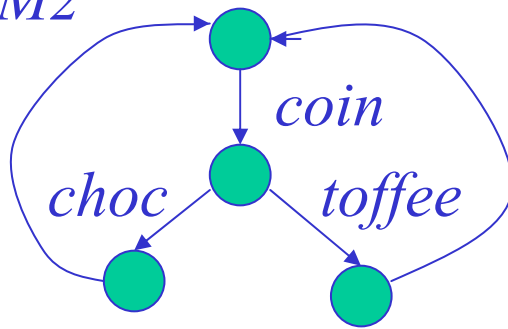


# Pictures

$VM2 = \mu X : \{coin, choc, toffee\} \cdot$

$coin \rightarrow (choc \rightarrow X \mid toffee \rightarrow X)$

$VM2$



Same process and different pictures  $\Rightarrow$  proofs of equality are difficult to conduct

# Pictures

- Another problem with pictures: processes with a very large or infinite number of states.
- Example : Rocket

$$\alpha\text{Rocket} = \{up, down, around\}$$

$$\text{Rocket} = R_0 = (up \rightarrow R_1 \mid around \rightarrow \text{Rocket})$$

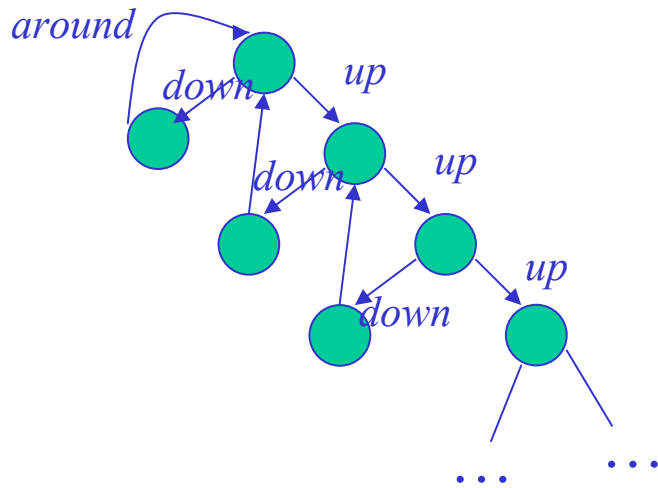
$$R_{n+1} = (up \rightarrow R_{n+2} \mid down \rightarrow R_n)$$

# Pictures: Exercise

1. Express Rocket as a tree structure.

# Pictures: Exercise

1. Express Rocket as a tree structure.



# Algebraic Laws of CSP

- There are many different ways of describing the same behaviour

$$(x \rightarrow P \mid y \rightarrow Q) = (y \rightarrow Q \mid x \rightarrow P)$$

- We must learn to recognize which expressions describe the same object and which not, as in arithmetic.

$$(x + y) = (y + x)$$

- CSP has a set of algebraic laws that defines its semantics and can be used to prove program equivalence.

# Some Laws of CSP

- L1: If  $A = B$  then  $(x:A \rightarrow P(x)) = (y:B \rightarrow P(y))$
- Consequences:
  - $STOP \neq (d \rightarrow P)$

Proof:

$$\begin{aligned} STOP &= (x:\{\} \rightarrow P) && \textit{by definition} \\ &\neq (x:\{d\} \rightarrow P) && \textit{because } \{\} \neq \{d\} \\ &= (d \rightarrow P) && \textit{by definition} \end{aligned}$$

# Some Laws of CSP

- L1: If  $A = B$  then  $(x:A \rightarrow P(x)) = (y:B \rightarrow P(y))$
- Consequences:
  - $(c \rightarrow P) \neq (d \rightarrow Q)$  if  $c \neq d$

Proof ?

# Some Laws of CSP

- L1: If  $A = B$  then  $(x:A \rightarrow P(x)) = (y:B \rightarrow P(y))$
- Consequences:
  - $(c \rightarrow P \mid d \rightarrow Q) = (d \rightarrow Q \mid c \rightarrow P)$

Proof:

Define  $R(x) = P$  if  $x=c$   
 $= Q$  if  $x=d$

$LHS = (x:\{c,d\} \rightarrow R(x))$       *by definition*  
 $= (x:\{d,c\} \rightarrow R(x))$       *because  $\{c,d\} = \{d,c\}$*   
 $= RHS$

# Some Laws of CSP: Exercise

Prove that

$$\begin{aligned} & \mu X : \{coin, choc, toffee\} \cdot coin \rightarrow (choc \rightarrow X \mid toffee \rightarrow X) \\ = & \mu X : \{coin, choc, toffee\} \cdot coin \rightarrow (toffee \rightarrow X \mid choc \rightarrow X) \end{aligned}$$

# Some Laws of CSP: Exercise

Prove that

$$\begin{aligned} & \mu X : \{coin, choc, toffee\} \bullet coin \rightarrow (choc \rightarrow X \mid toffee \rightarrow X) \\ = & \mu X : \{coin, choc, toffee\} \bullet coin \rightarrow (toffee \rightarrow X \mid choc \rightarrow X) \end{aligned}$$

Proof:

*Immediate consequence of the previous result.*

# Internal (nondeterministic) choice

- If  $P$  and  $Q$  are processes then

$$P \sqcap Q$$

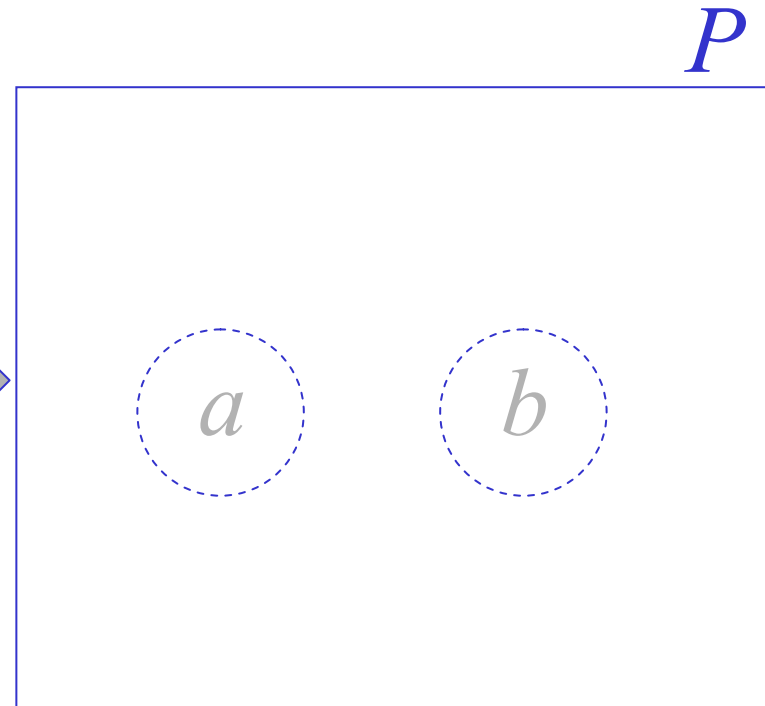
denotes the process which is either  $P$  or  $Q$ ; the choice is arbitrary, without the influence of the environment.

- The alphabet of  $P$  and  $Q$  is the same.
- Internal choice **is not a combinator** an implementor would wish to use in combining implementations. We are forced to consider it since nondeterminism arises from abstraction.

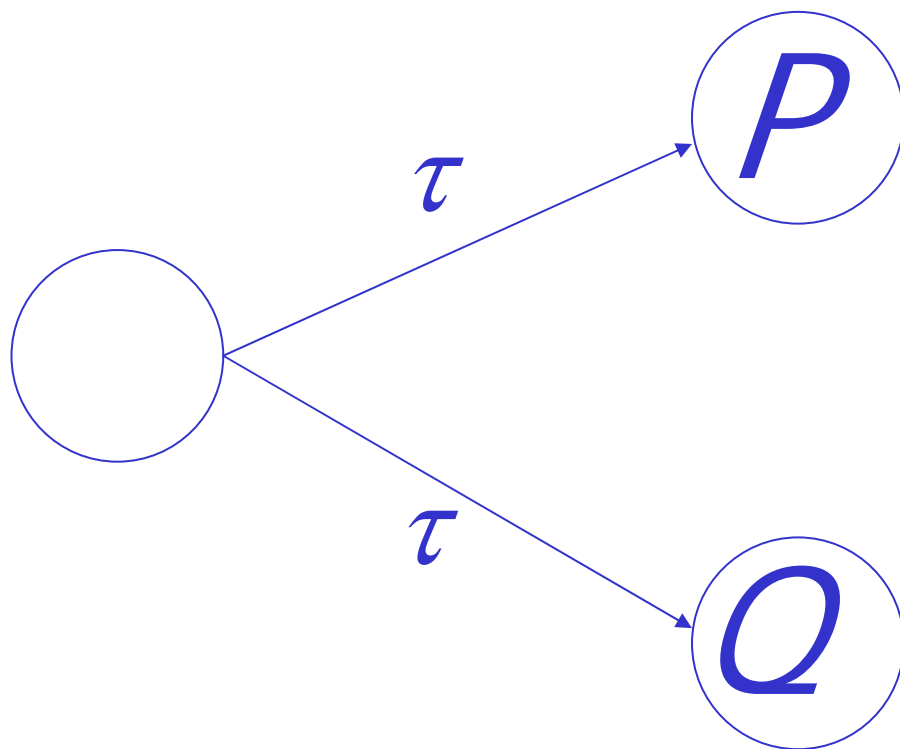
# Internal Choice

Notation:  $\sqcap$  or  $|\sim|$

$$P = a \rightarrow Q$$
$$|\sim| b \rightarrow R$$



# Internal Choice ( $P \sim Q$ )



# Internal Choice

- How internal choice arises?

Recall that in sequential programming, abstracting boolean variable  $x$  in program

*if  $x = true$  then  $P$  else  $Q$*

gives program

$\llbracket \text{var } x \bullet \text{if } x = true \text{ then } P \text{ else } Q \rrbracket$

which is equivalent to a nondeterministic choice between  $P$  and  $Q$ :

$P \sqcap Q.$

# Internal Choice

- Examples:

Recall the change-giving machine

$$\alpha Chh = \{in5, out1, out2\}$$

$$Chh = (in5 \rightarrow (out2 \rightarrow out1 \rightarrow out1 \rightarrow out1 \rightarrow Chh \\ | out1 \rightarrow out2 \rightarrow out2 \rightarrow Chh))$$

For comparison consider

$$Ch1 = in5 \rightarrow Ones$$

$$Ones = out2 \rightarrow out1 \rightarrow out1 \rightarrow out1 \rightarrow Ch1$$

$$Ch2 = in5 \rightarrow Twos$$

$$Twos = out1 \rightarrow out2 \rightarrow out2 \rightarrow Ch2$$

$$Chh = Ch1 \sqcap Ch2$$

Are they equivalent?

# Internal Choice

- Consequences of internal choice
  - 1. The appearance of internal choice in a design means that in reasoning about the design the designer must encompass all possible implementations; **bad**.
  - 2. The appearance of internal choice in a design allows the implementor a choice in finding an implementation; **good**. Process  $P \sqcap Q$  **may be implemented by** either  $P$  or  $Q$  (a choice for the implementor).

# Internal Choice: Laws

- Internal choice is set-like (idempotent, commutative and associative)

$$P \sqcap P = P$$

$$P \sqcap Q = Q \sqcap P$$

$$P \sqcap (Q \sqcap R) = (P \sqcap Q) \sqcap R.$$

- A choice that first does  $a$  and then makes a choice is indistinguishable from one which first makes the choice and then does  $a$ .

$$a \rightarrow (P \sqcap Q) = (a \rightarrow P) \sqcap (a \rightarrow Q)$$

- General choice distributes through internal choice.

$$?x : A \rightarrow (P(x) \sqcap Q(x)) = (?x : A \rightarrow P(x)) \sqcap (?x : A \rightarrow Q(x))$$

# Internal Choice: Laws

- Attention: Recursion **does not** distribute through internal choice. Let  $a \neq b$  and

$$P = \mu X. ((a \rightarrow X) \sqcap (b \rightarrow X))$$

$$Q = (\mu X. (a \rightarrow X)) \sqcap (\mu X. (b \rightarrow X))$$

then  $P \neq Q$ .

Why?

## Internal Choice: Exercise

- A bus company guarantees to provide buses between  $A$  and  $B$ , but does not guarantee any particular route. There are two routes, 100 and 200, and the passenger is happy to accept either. Model the service offered by the bus company.

# Internal Choice: Exercise

- A bus company guarantees to provide buses between  $A$  and  $B$ , but does not guarantee any particular route. There are two routes, 100 and 200, and the passenger is happy to accept either. Model the service offered by the bus company.

$$BUS = BUS\_100 \sqcap BUS\_200$$

# External choice

- Process  $(a \rightarrow A \mid b \rightarrow B)$  is also written

$$(a \rightarrow A) \square (b \rightarrow B)$$

in which now  $a$  and  $b$  need not be distinct.

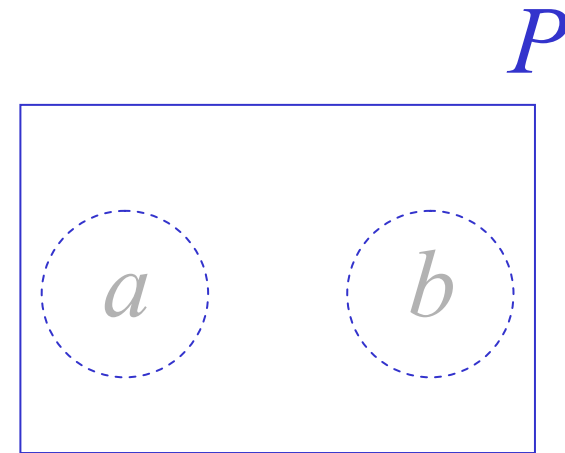
- If  $P$  and  $Q$  are processes then

$$P \square Q$$

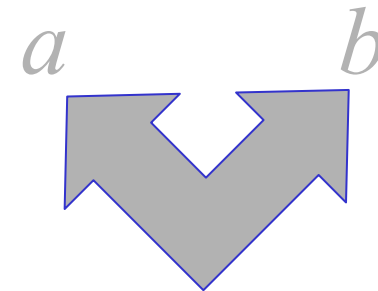
denotes their external choice whose initial event is determined by the environment: if it is an event of just  $P$  then  $P \square Q$  behaves like  $P$ ; if it is an event of just  $Q$  then it behaves like  $Q$ ; if it is an event common to both  $P$  and  $Q$  then its behaviour is the nondeterministic choice between  $P$  and  $Q$ .

# External Choice

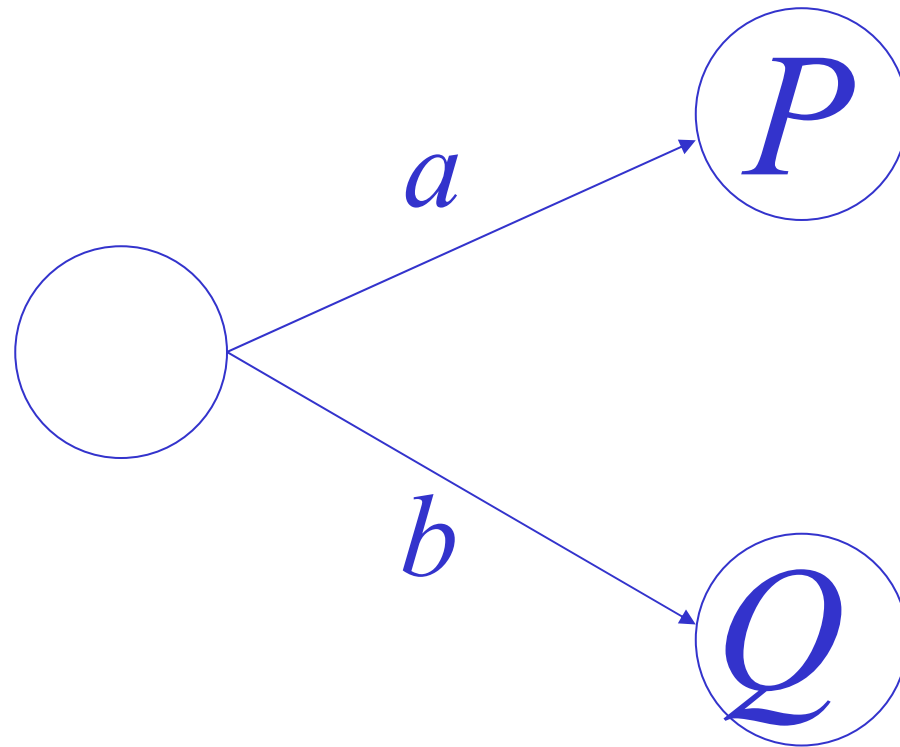
$$P = a \rightarrow Q$$
$$[ ] b \rightarrow R$$



**OR**



# External Choice ( $P \ [ \ ] \ Q$ )



Provided that  $a$  is a starting event of  $P$  and  $b$  a starting event of  $Q$

# External choice

- The behaviour of  $P \square Q$  is thus:
  - if the environment's first action is not possible for either  $P$  or  $Q$  then it does not occur;
  - if the environment's first action is possible only for  $P$  then  $P$  is selected;
  - if the environment's first action is possible only for  $Q$  then  $Q$  is selected;
  - if the environment's first action is possible for both  $P$  and  $Q$  then  $P \sqcap Q$  is selected.

- Thus,

$$\begin{aligned}(a \rightarrow P) \square (b \rightarrow Q) &= (a \rightarrow P \mid b \rightarrow Q), a \neq b, \\ &= (a \rightarrow P) \sqcap (a \rightarrow Q), a = b.\end{aligned}$$

# External choice

- Example:
  - A vending machine can offer chocolate and toffee. The user will select the sweet he want and after pay for the appropriate amount.

# External choice

- Example:
  - A vending machine can offer chocolate and toffee. The user will select the sweet he want, after paying for the appropriate amount.

$$VM = \mu X : \{coin10, coin50, choc, toffee\} \bullet \\ (coin10 \rightarrow toffee \rightarrow X [] coin50 \rightarrow choc \rightarrow X)$$

# Comparing internal and external choices

- Observe that the process

$$(a \rightarrow A) \sqcap (b \rightarrow B)$$

offer its environment a choice between  $a$  and  $b$  initially.  
However the process

$$(a \rightarrow A) \sqcap (b \rightarrow B)$$

offers either  $a$  or  $b$  and internally decides which, without the influence of the environment.

- The latter may lead to deadlock if the environment offers only  $a$  whilst the former cannot.

# External Choice: Laws

- External is set-like (idempotent, commutative and associative)

$$P \sqcap P = P$$

$$P \sqcap Q = Q \sqcap P$$

$$P \sqcap (Q \sqcap R) = (P \sqcap Q) \sqcap R.$$

.

- STOP is a unit:  $P \sqcap STOP = P$ .

- General law:

$$(?x:A \rightarrow P(x)) \sqcap (?y:B \rightarrow Q(y))$$

=

$$?z: A \cup B \rightarrow P(z), \quad z \in A \setminus B$$

$$Q(z), \quad z \in B \setminus A$$

$$P(z) \sqcap Q(z), \quad z \in A \cap B$$

# External choice: Laws

- External choice distributes through internal choice and vice-versa.

$$P \square (Q \sqcap R) = (P \square Q) \sqcap (P \square R)$$

$$P \sqcap (Q \square R) = (P \sqcap Q) \square (P \sqcap R)$$

- Nevertheless,

$$a \rightarrow (P \square Q) \neq (a \rightarrow P) \square (a \rightarrow Q)$$

Why?

# Input and External Choice

Let  $c$  be a channel of type  $\{e_0, \dots, e_N\}$

$c?x \rightarrow P$



$c.e_0 \rightarrow P[e_0/x]$

[ ]

...

[ ]

$c.e_N \rightarrow P[e_N/x]$

# External Choice: Exercise

1. Present a CSP process which captures the first menu of an ATM machine.

# External Choice: Exercise

1. Present a CSP process which captures the first menu of an ATM machine.

*Menu = withdraw -> P*

*[] balance -> Q*

*[] deposit -> R ...*

# Conditional Choice

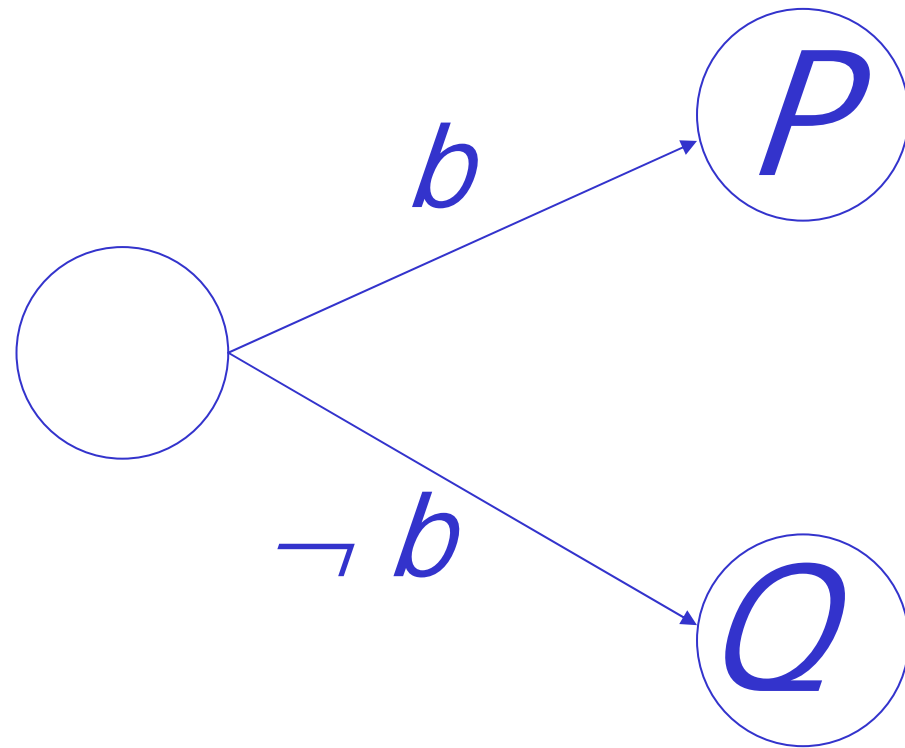
- If  $b$  is a predicate and  $P$  and  $Q$  are processes then

$$P \triangleleft b \triangleright Q$$

denotes the process which behaves like  $P$  if  $b$  holds and like  $Q$  otherwise. It is read  $P$  if  $b$  else  $Q$ .

- Predicate  $b$  is on the appropriate ‘state space’. Its evaluation is not an event of the conditional process.
- The process  $b \ \& \ P$  (guard) is a shorthand of  $P$  if  $b$  else STOP

# Conditional Choice



# Conditional Choice: Laws

- The operator  $\triangleleft b \triangleright$  is idempotent, associative and distributes through operator  $\triangleleft c \triangleright$ .

- Conditional choice satisfies the propositional laws

$$P \triangleleft true \triangleright Q = P$$

$$P \triangleleft b \triangleright Q = Q \triangleleft \neg b \triangleright P \dots$$

- It distributes through both internal and external choice in each argument

$$(P \sqcap Q) \triangleleft b \triangleright R = (P \triangleleft b \triangleright R) \sqcap (Q \triangleleft b \triangleright R)$$

$$P \triangleleft b \triangleright (Q \sqcap R) = (P \triangleleft b \triangleright Q) \sqcap (P \triangleleft b \triangleright R)$$

$$(P \sqcup Q) \triangleleft b \triangleright R = (P \triangleleft b \triangleright R) \sqcup (Q \triangleleft b \triangleright R)$$

$$P \triangleleft b \triangleright (Q \sqcup R) = (P \triangleleft b \triangleright Q) \sqcup (P \triangleleft b \triangleright R)$$

# Conditional Choice: Laws

- General case:

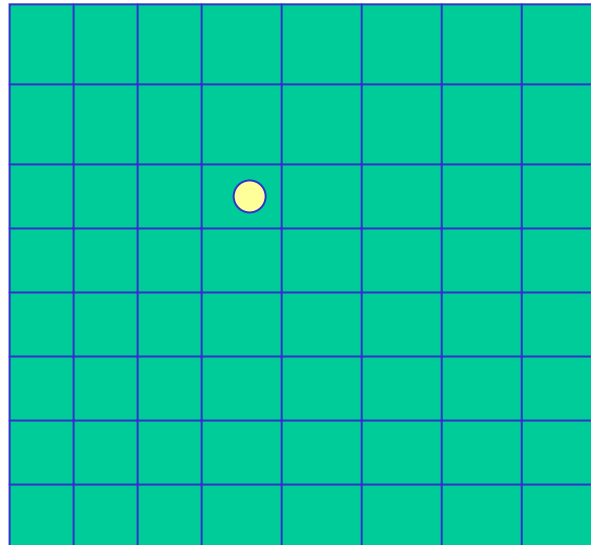
$$?x:A \rightarrow (P(x) \triangleleft b \triangleright Q(x)) = (?x:A \rightarrow P(x)) \triangleleft b \triangleright (?x:A \rightarrow Q(x))$$

- Conditional choice is typically used after an input event. The general law for external choice can be expressed using conditional choice:

$$\begin{aligned} & (?x:A \rightarrow P(x)) \sqcap (?y:B \rightarrow Q(y)) \\ = & \\ & ?z: A \cup B \rightarrow (P(z) \sqcap Q(z)) \\ & \triangleleft z \in A \cap B \triangleright \\ & (P(z) \triangleleft z \in A \triangleright Q(z)) \end{aligned}$$

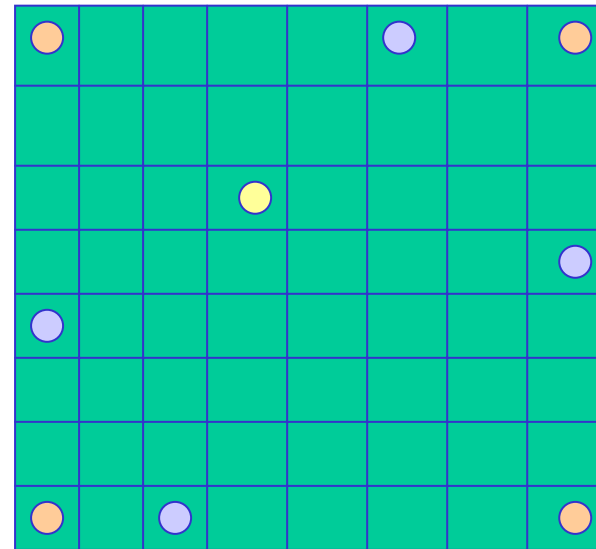
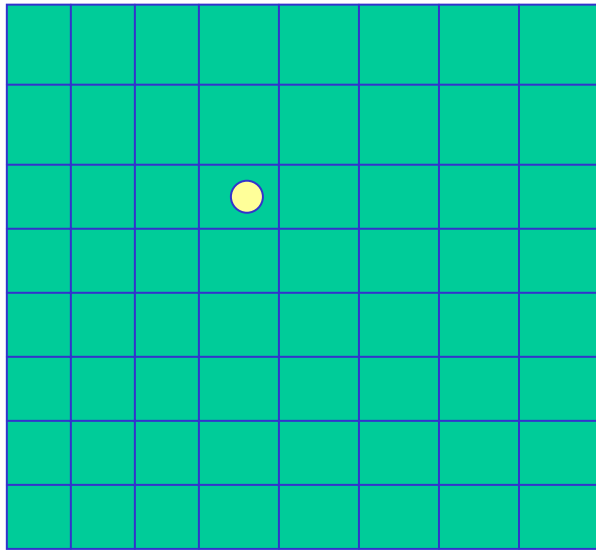
# Conditional Choice: Example

- Conditional choice can thus be used to simplify process definitions . Consider the board below. How to describe the possible movements of the yellow ball?



# Conditional Choice: Example

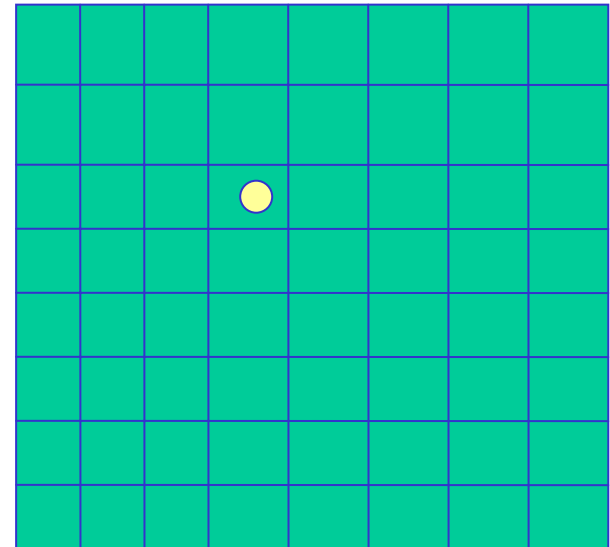
- We need to consider nine cases.



# Conditional Choice: Example

- Old approach. First for an interior state  
 $0 < x, y < 7$ :

$$C(x,y) = \begin{array}{l} (up \rightarrow C(x, y + 1) \\ | down \rightarrow C(x, y-1) \\ | left \rightarrow C(x-1, y) \\ | right \rightarrow C(x+1, y) \end{array}$$



# Conditional Choice: Example

- Then eight more cases !!!!

Cases:

$$x = y = 0,$$

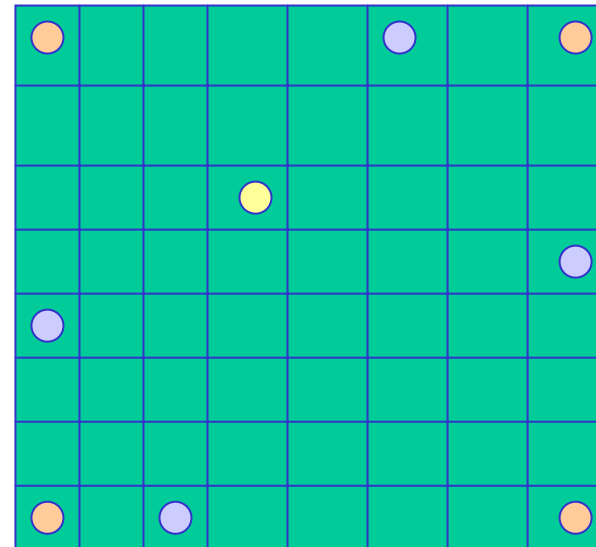
$$x = 0 \text{ and } y = 8,$$

$$x = 8 \text{ and } y = 0,$$

$$x=8 \text{ and } y = 8,$$

$$x = 0 \text{ and } 0 < y < 7$$

...



# Conditional Choice: Example

- Using conditional choice, we have a **single case**:

$$C(x,y) = ( \quad \textit{up} \rightarrow C(x, y + 1) \triangleleft y < 7 \triangleright STOP$$
$$\quad \quad \quad [] \textit{down} \rightarrow C(x, y-1) \triangleleft 0 < y \triangleright STOP$$
$$\quad \quad \quad [] \textit{left} \rightarrow C(x-1, y) \triangleleft 0 < x \triangleright STOP$$
$$\quad \quad \quad [] \textit{right} \rightarrow C(x+1,y) \triangleleft x < 7 \triangleright STOP \}$$

# Conditional Choice: Exercise

1. Write a CSP process that accepts a value and if the value is positive outputs its double and if it is negative does nothing.

# Conditional Choice: Exercise

1. Write a CSP process that infinitely accepts a value. If the value is positive outputs its double, otherwise does nothing.

$$P = in ? x \rightarrow ((out ! 2 * x \rightarrow P) \triangleleft x > 0 \triangleright P)$$